

Quantum networks  
 $U: (\mathbb{C}^2)^{\otimes n} \rightarrow (\mathbb{C}^2)^{\otimes n}$  unitary.  
 universal gates? yes, today

Examples for 1-qbit gates

$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$  Hadamard gate

$\sigma_1, \sigma_2, \sigma_3$   $\sigma_1 = \text{NOT}$   
 (in the standard basis)  
 $\sigma_1 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$   $\sigma_1 \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$   $\sigma_1 \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$

$S(\varphi) = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\varphi} \end{pmatrix}$

Def. (i) for  $x_0, \dots, x_n$  let  $|x_0 \dots x_n\rangle = |x_0\rangle \otimes \dots \otimes |x_n\rangle$   
 $x_i \in \{0, 1\}$

$x = x_0 + 2x_1 + 4x_2 + \dots + 2^n x_n \in \mathbb{N}$   
 $|x\rangle_n := |x_0 \dots x_n\rangle_n$

[Ex:  $|0\rangle_2 = |00\rangle_2 = |0\rangle \otimes |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$   
 $|1\rangle_2 = |01\rangle_2 = |0\rangle \otimes |1\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$   
 $|2\rangle_2 = |10\rangle_2 = |1\rangle \otimes |0\rangle = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$   
 $|3\rangle_2 = |11\rangle_2 = |1\rangle \otimes |1\rangle = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$

(ii) for  $U = \begin{pmatrix} u_{00} & u_{01} \\ u_{10} & u_{11} \end{pmatrix} \in U(2)$  and  $m \in \mathbb{N}_0$

$\Lambda_0(U) := U$

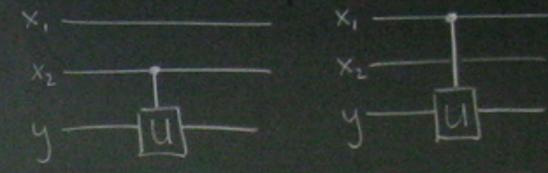
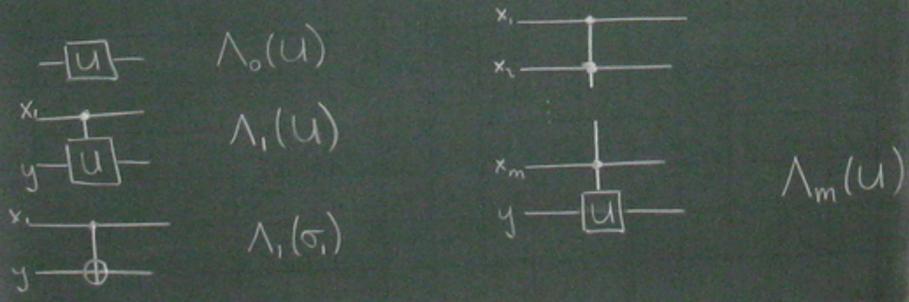
$\Lambda_n(U) |x_1 \dots x_m y\rangle_{n+1} = \begin{cases} |x_1 \dots x_m y\rangle_{n+1} & \text{if } x_1 \dots x_m = 0 \\ |x_1\rangle \otimes \dots \otimes |x_m\rangle \otimes U|y\rangle & \text{if } x_1 \dots x_m = 1 \\ (\Leftrightarrow x_1 = x_2 = \dots = x_m = 1) \end{cases}$

$\Lambda_m(U)$  controlled U-gate with m control bits

(iii) The Toffoli gate is  $\Lambda_1(\sigma_1)$   
 (controlled-NOT gate)

graphical notation

$\begin{matrix} 1 \\ 2 \\ \vdots \\ n \end{matrix} \begin{matrix} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{matrix} \leftrightarrow |x_1\rangle \otimes \dots \otimes |x_n\rangle$



matrix rep. of  $\Lambda_1(U)$   $\Lambda_1(U) |0\rangle \otimes |y\rangle = |0\rangle \otimes |y\rangle$   
 $\Lambda_1(U) |1\rangle \otimes |y\rangle = |1\rangle \otimes U|y\rangle$

$|0\rangle \otimes |y\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} 1 \cdot y_1 \\ 0 \cdot y_1 \\ 1 \cdot y_2 \\ 0 \cdot y_2 \end{pmatrix} = \begin{pmatrix} y_1 \\ 0 \\ y_2 \\ 0 \end{pmatrix}$

$|1\rangle \otimes |y\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} 0 \cdot y_1 \\ 1 \cdot y_1 \\ 0 \cdot y_2 \\ 1 \cdot y_2 \end{pmatrix} = \begin{pmatrix} 0 \\ y_1 \\ 0 \\ y_2 \end{pmatrix}$

$\Rightarrow \Lambda_1(U) = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & u_{00} & u_{01} \\ 0 & 0 & u_{10} & u_{11} \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & U \end{pmatrix}$

Theorem. Every unitary  $U$  on  $\mathcal{H}_n = (\mathbb{C}^2)^{\otimes n}$   
 can be written as a product of  
 1-qbit operators  $\Lambda_0(V)$ ,  $V \in U(2)$   
 and the 2-qbit  $\Lambda_1(\sigma_1) = T$

Proof in many little steps (sketched here)

Lemma 1.  $U \in U(2) \rightarrow \exists \alpha, \beta, \delta, \theta \in \mathbb{R}$  such that

$$U = e^{i\alpha} e^{i\frac{\beta}{2}\sigma_1} e^{i\frac{\delta}{2}\sigma_2} e^{i\frac{\theta}{2}\sigma_3}$$

$U \in SU(2) \Rightarrow \delta = 0$

$$U = e^{i\alpha} \begin{pmatrix} e^{i\theta} & 0 \\ 0 & e^{-i\theta} \end{pmatrix} \begin{pmatrix} \cos \frac{\beta}{2} & i \sin \frac{\beta}{2} \\ i \sin \frac{\beta}{2} & \cos \frac{\beta}{2} \end{pmatrix} \begin{pmatrix} e^{i\alpha} & 0 \\ 0 & e^{-i\alpha} \end{pmatrix}$$

Proof Exercise!

Lemma 2.  $W \in SU(2) \rightarrow \exists A, B, C \in SU(2)$   
 $ABC = 1$  and  $A\sigma_1 B\sigma_1 C = W$

Proof. For  $k=1,2,3$  let  $R_k(\alpha) = e^{i\frac{\alpha}{2}\sigma_k}$

Lemma 1  $\Rightarrow \exists \alpha, \beta, \theta: W = R_3(\alpha)R_2(\theta)R_3(\beta)$

Let  $A = R_3(\alpha)R_2(\frac{\theta}{2})$

$$B = R_2(-\frac{\theta}{2})R_3(-\frac{\alpha+\beta}{2}), C = R_3(\frac{\beta-\alpha}{2})$$

$\Rightarrow ABC = 1$

$$A\sigma_1 B\sigma_1 C = R_3(\alpha)R_2(\frac{\theta}{2})\sigma_1 R_2(-\frac{\theta}{2})R_3(-\frac{\alpha+\beta}{2})\sigma_1 R_3(\frac{\beta-\alpha}{2})$$

$$= R_3(\alpha)R_2(\theta)R_3(\alpha) = W$$

Lemma 3  $\delta \in \mathbb{R} \rightarrow \Lambda_1(e^{i\delta} 1) = E \otimes 1$

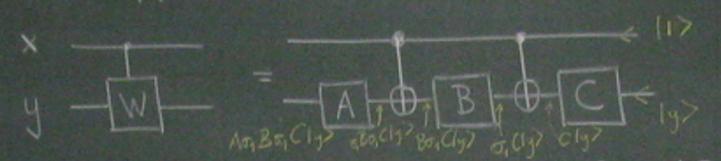
where  $E = R_3(-\delta) e^{i\frac{\delta}{2}\sigma_3}$  is unitary.



Proof Exercise

Lemma 4  $W \in SU(2)$  Then

$$\Lambda_1(W) = (1 \otimes A) T (1 \otimes B) T (1 \otimes C)$$

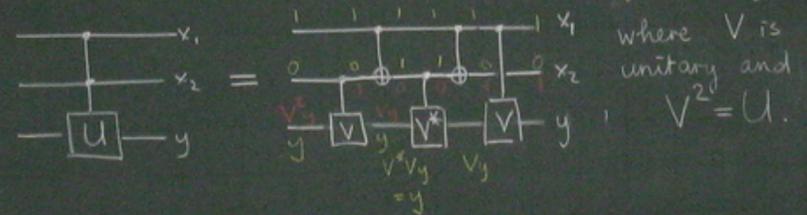


Proof:  $x=0: A \cdot 1 \cdot B \cdot 1 \cdot C = ABC = 1$  acting on  $|y\rangle$ .

$x=1: A\sigma_1 B\sigma_1 C = W$  acting on  $|y\rangle$ .

Corollary:  $W \in U(2) \Rightarrow$  apply Lemma 3 to get:  
 $\Lambda_1(W)$  is a product of 2 T's and 4 1-qbit gates.

Lemma 5.  $U \in U(2)$ . Then  $\Lambda_2(U)$  is given by



Corollary 2  $\Lambda_2(\sigma_1)$  is implementable by T and single-qbit gates ( $U = \sigma_1$ )

recall  $\Lambda_2(\sigma_1)$  is the 3 bit Toffoli gate used to generate classical networks.  
 A quantum network can perform any classical computation

In particular, any permutation of the basis states can be done using T and  $\sqrt{\sigma_1}$

Lemma 6.  $n \geq 3, U \in U(2)$ . Then  $\Lambda_{n+1}(U)$  can be constructed as an n-bit network of T's and  $\Lambda_1(V)$  or  $\Lambda_1(V^*)$  with V unitary and  $V^{(2^{n-2})} = U$ .

Lemma 7  
 think of  $d=2^n$

$U \in U(d) \rightarrow U$  is a product of  $\leq 2d^2$  unitary matrices, each of which acts only in a two-dimensional subspace of  $\mathbb{C}^d$ .

Proof

$\forall v \in \mathbb{C}^d, |v|=1 \Rightarrow \exists d-1$  such unitaries that transform  $v$  to  $\begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$

$$A(v_1, v_2) = \frac{1}{\sqrt{|v_1|^2 + |v_2|^2}} \begin{pmatrix} \bar{v}_1 & \bar{v}_2 \\ v_2 & -v_1 \end{pmatrix} \in U(2)$$

$$A(v_1, v_2) \cdot \begin{pmatrix} v_1 \\ v_2 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \sqrt{|v_1|^2 + |v_2|^2}$$

$$v \in \mathbb{C}^d \begin{pmatrix} v_1 \\ v_2 \\ \vdots \\ v_d \end{pmatrix} \xrightarrow{A} \begin{pmatrix} \sqrt{|v_1|^2 + |v_2|^2} \\ 0 \\ \vdots \\ v_d \end{pmatrix} \rightarrow \dots \rightarrow \begin{pmatrix} \sqrt{|v_1|^2 + |v_2|^2} \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$